

Recall A game involves

- $n$  players
- a set  $S_i$  of strategies for player  $i$
- a payoff function

$$v_i : S_1 \times S_2 \times \dots \times S_n \rightarrow \mathbb{R}$$

for player  $i$ ,  $1 \leq i \leq n$ .

Example  $n=2$

$$S_1 = \{H, T\}, \quad S_2 = \{H, T\}$$

$$v_1(H, H) = 1 \quad v_1(H, T) = -1$$

$$v_2(H, H) = -1 \quad v_2(H, T) = 1$$

$$v_1(T, H) = -1 \quad v_1(T, T) = 1$$

$$v_2(T, H) = 1 \quad v_2(T, T) = -1$$

Recall A mixed strategy is a choice of probabilities

$$p_{i,s}$$

for  $1 \leq i \leq n$ ,  $s \in S_i$  satisfying

$$p_{i,s} \geq 0, \quad \sum_{s \in S_i} p_{i,s} = 1.$$

Here  $p_{i,s}$  is the probability that player  $i$  plays strategy  $s \in S_i$ .

Notation Suppose  $S_i = \{s_1, s_2, \dots, s_k\}$ , and set

$$\underline{p}_i = (p_{i,s_1}, p_{i,s_2}, \dots, p_{i,s_k}) \in \mathbb{R}^{|S_i|}$$

Define the expected payoff for player  $i$  to be the function

$$E_i: \mathbb{R}^{|S_1|+|S_2|+\dots+|S_n|} \rightarrow \mathbb{R}$$

by

$$E_i(\underline{p}_1, \underline{p}_2, \dots, \underline{p}_n) = E(v_i)$$

$$= \sum_{\substack{x_1 \in S_1 \\ x_2 \in S_2 \\ \vdots \\ x_n \in S_n}} p_{1,x_1} p_{2,x_2} \dots p_{n,x_n} v_i(x_1, x_2, \dots, x_n)$$

Example For our 2-player game

$$S_1 = \{H, T\}, S_2 = \{H, T\}$$

$$\Sigma_i(\underline{P}_1, \underline{P}_2) =$$

$$P_{1H} P_{2H} v_i(H, H) + P_{1T} P_{2H} v_i(T, H)$$

$$+ P_{1H} P_{2T} v_i(H, T) + P_{1T} P_{2T} v_i(T, T)$$

$$\Sigma_1(P_{1H}, P_{1T}, P_{2H}, P_{2T}) = P_{1H} P_{2H} - P_{1T} P_{2H}$$

$$- P_{1H} P_{2T} + P_{1T} P_{2T}$$

Theorem (Nash)

Example In our 2-player game  
an example of a mixed Nash  
equilibrium is the strategy

$$P_{1H} = \frac{1}{2} \quad P_{1T} = \frac{1}{2} \quad P_{2H} = \frac{1}{2} \quad P_{2T} = \frac{1}{2}$$



Theorem (Nash) For any game with finitely many players and finite strategy sets there exists at least one mixed Nash equilibrium.

Outline proof

Consider

$$C = \{ (p_1, p_2, \dots, p_n) \} \subseteq \mathbb{R}^{|S_1| + |S_2| + \dots + |S_n|}$$

where  $\underline{p}_i \in \mathbb{R}^{|S_i|}$  is the probability distribution for player  $i$ ,  $1 \leq i \leq n$ .

Now  $C$  is closed, bounded and

convex, since  $p_{i,s} \geq 0$ ,  $\sum_{s \in S_i} p_{i,s} = 1$ .

Thus any continuous function

$f: C \rightarrow C$  has at least one

fixed point (Brouwer's Theorem).

For a given  $(p_1, p_2, \dots, p_n) \in C$   
define  $q_i \in R^{I_i}$  to be the  
probability distribution  $q$  that  
maximizes

$$\max_q \sum_i (p_1, p_2, \dots, p_{i-1}, q, p_{i+1}, \dots, p_n) \quad (*)$$

Now define  $f: C \rightarrow C$  by

$$f(p_1, p_2, \dots, p_n) = (q_1, q_2, \dots, q_n)$$

This function  $f$  has a fixed  
point. But a fixed point is  
a Nash equilibrium. □

Slight problem: The quantity  $(*)$   
might achieve its max. value  
at more than one value of  $q$ .  
So our  $f$  is not well-defined!

To solve this problem replace (\*)  
by

$$\max_{\underline{q}} \left( \sum_i (\underline{p}_1, \dots, \underline{p}_{i-1}, \underline{q}, \underline{p}_{i+1}, \dots, \underline{p}_n) - \|\underline{p}_i - \underline{q}\|^2 \right) \quad (**)$$