

# RSA Public Key Cryptography

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(Rivest, Shamir, Adleman 1977)

Suppose

- $N$  letter alphabet (e.g.  $N=26$ )
- $k$ -letter plaintext message units
- $l$ -letter ciphertext message units

Plaintext  
message  
units



Integers

$$0 \leq i \leq N^k$$

Ciphertext  
message  
units



Integers

$$0 \leq c \leq N^l$$

## Crypto system

- Each user chooses two distinct random prime numbers  $P, Q$  (of around 100 digits each to be safe with current technology),

HELLO THERE

H E L L

HE LL OT HE RE

HEL LOT HER E--

- Choose an integer  $e$  with  $\text{gcd}(e, p-1) = 1 = \text{gcd}(e, q-1)$  (2)

- Each user computes

$$n = pq$$

and publishes the enciphering key

$$K_E = (n, e).$$

- Each user computes (using the Euclidean algorithm)

$$d = e^{-1} \text{ mod } \phi(n)$$

where  $\phi(n) = (p-1)(q-1)$

The deciphering key

$$K_D = (n, d)$$

is kept secret.



- The enciphering function is 3
- $$f_{(n,e)} : \mathbb{Z}_n \rightarrow \mathbb{Z}_n, x \mapsto x^e$$

Proposition

$$(x^e)^d = x \pmod{n}$$

- The deciphering function is
- $$f_{(n,d)} : \mathbb{Z}_n \rightarrow \mathbb{Z}_n, x \mapsto x^d.$$

## Example of RSA cryptosystem

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26-letter alphabet  $A=0, B=1, \dots, Z=25$

$k=3$  : 3-letter plaintext units

$l=4$  : 4-letter ciphertext units

I want to send Alice the message

YES

Here published public key is

$$K_E^{\text{Alice}} = (n, e)$$

$$= (46927, 39423)$$

$$\text{YES} \leftrightarrow 24 \cdot 26^2 + 4 \cdot 26 + 18$$

$$= 16346$$

$$f_{(n,e)}^{\text{Alice}}(\text{YES}) = 16346^{39423} \bmod 46927$$

$$\equiv 21166$$

$$21166 = 1 \cdot 26^3 + 5 \cdot 26^2 + 8 \cdot 26 + 2$$

(5)

$$= \text{BFIC}$$

to send the enciphered message  
BFIC.

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It is believed that the  
computation of  $d$  necessitates  
the factorisation of  $n$  into

$$n = p q$$

It is believed that (with  
current methods) the factorization  
would take a prohibitively long  
time.